## Selecting the Most Conflicting Pair of Candidates

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## setting the stage

$$
\begin{aligned}
& \text { 田〉0>畟〉良〉为 }
\end{aligned}
$$


for what?


Learning from preferences

## engagement

## creativity

deliberation


## selecting: state of the art

## Most

conflicting candidates


## Current tools

 insufficient!preference insights: state of the art
＊

## Single－candidate measures



Full－election measures

$$
\begin{aligned}
& \text { カ〉 }
\end{aligned}
$$

## Current tools

 insufficient！
## must-have properties




## Unanimity



## Conflict consistency and unanimity are contradicting each other!

nice-to-have properties

$V^{a>b}$ voters preferring $a$ to $b$
$v(a b)$ "directed" positions difference between $a$ and $b$


Matching-domination of pairs (informally)

Pair $\{A, B\}$ dominates pair $\{C, D\}$ if voters can be matched such that for each matched pair the conflict between $A$ and $B$ is at least that between $C$ and $D$; with these inequality being strong for at least one pair. Each matched pair of voters has the same preference towards $\{A, B\}$ and $\{C, D\}$.

|  | $\begin{aligned} & A>B \\ & v(A B) \end{aligned}$ | $\underset{v(C D)}{C>D}$ | $A>B$ | C>D |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $A>B>C>D$ | 1 | 1 |  | $\geq 2$ |  |
| $A>C>D>B$ | 3 | 1 |  | $\geq 1$ |  |
| $A>C>B>D$ | 2 | 2 |  | $\geq 1$ |  |
| $D>C>A>B$ | 1 | -1 |  | $\geq 1$ | $A>B$ is dominating $D>C$ |
|  | $\begin{aligned} & B>A \\ & v(B A) \end{aligned}$ | $\begin{aligned} & \mathrm{D}>\mathrm{C} \\ & v(D C) \end{aligned}$ | $B>A$ | D $>C$ |  |
| $B>D>A>C$ | 2 | 2 | 3 | $\geq 2$ |  |
| $B>C>D>A$ | 3 | -1 |  | $\geq 1$ |  |

Matching domination
Matching-dominated pairs are never selected!
nobody's perfect

Theorem: Conflict consistency, matching domination, and conflict monotonicity are incompatible.

Proof:
$a>b>c>d$
$b>a>d>c$
Only (a,b) or (c,d) can win

## Conflict Consistency

Assume ( $\mathrm{a}, \mathrm{b}$ ) is wins
$a>b>c>d$
$b>d>c>a$
(a,b) should still win
$a>b: \quad(1,-3)$
$a>d: \quad(-3,2)$
(a,d) dominates (a,b),
Conflict Monotonicity
thus (a,b) cannot win

## getting the most conflicting pair

Conflict between two voters $\operatorname{conf}_{v, v^{\prime}}^{\circ}(a, b)= \begin{cases}0 & \text { if } v(a b) \cdot v^{\prime}(a b)>0 \\ |v(a b)| \circ\left|v^{\prime}(b a)\right| & \text { otherwise }\end{cases}$ $+$
de

## Max Sum Conflict

$\operatorname{MaxSum}(P)=\underset{a, b \in C}{\operatorname{argmax}} \sum_{v, v^{\prime} \in V} \operatorname{conf}^{\dagger}(a, b)$

## Max Nash Conflict

$\operatorname{MaxNash}(P)=\underset{a, b \in C}{\operatorname{argmax}} \sum_{v, v^{\prime} \in V} \operatorname{conf}^{\mathrm{x}}(a, b)$

$$
\begin{aligned}
& \operatorname{conf}^{+}(\mathbf{2}, 0)=0 \\
& \operatorname{conf}^{\times} \text {(田触) }=4 \cdot 4=16
\end{aligned}
$$

$$
\operatorname{nonconf}(a, b)=\min \left(\sum_{v \in V^{a>b}} v(a b), \sum_{v \in V^{b>a}} v(b a)\right)
$$



$$
\begin{aligned}
& \operatorname{nonconf}(\text { 田, })=\min (4,4)=4 \\
& \operatorname{nonconf}(\%, 0)=\min (2,0)=0 \\
& \operatorname{nonconf}(O, \boxplus)=\min (1,3)=1
\end{aligned}
$$

## Max Swap

$\operatorname{MaxSwap}(P)=\operatorname{argmax} \operatorname{nonconf}(a, b)$

$$
a, b \in C
$$

## Understanding <br> the ???

# Understanding <br> the conflictual <br> voting rules <br> (axiomatically) 



Axiomatic properties of conflictual rules.

# Understanding the conflictual voting rules 

(quantitatively)

Who is the most conflicting, $\{A, B\}$ or $\{X, Y\}$ ?

| Y | $A>X>\ldots>Y>\ldots>B$ |
| :---: | :---: |
| $B>X>\ldots>Y>A$ | $B>X>\ldots \ggg \ggg 1$ |
| $>Y>\ldots>X>A$ | $B>Y>\ldots$ |
| $B>Y>\ldots>X>A$ | > Y > ... > X > ... |

## Discrepancy

$$
\beta(a, b)=\frac{1}{n(m-1)} \sum_{v \in V} v(a b)
$$

Max Sum $\operatorname{conf}^{+}(a, b)=C_{2} \alpha(2-\alpha) \beta$
Max Nash $\operatorname{conf}^{\times}(a, b)=C_{1} \alpha(2-\alpha) \beta^{2}$ Max Swap nonconf $(a, b)=C_{3} \alpha \beta$

|  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| Reverse Stability | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ |
| Conflict Consistency | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ |
| Conflict Monotonicity | $x$ | $x$ | $x$ | $x$ |
| Antagonization Consistency | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ |
| Matching Domination | $\checkmark$ | $\checkmark$ | $\checkmark$ | $x$ |

Axiomatic properties of conflictual rules.

## Understanding the conflictual voting rules

(experimentally)


Politics

## X axis: partitioning ratio $\alpha$

the higher the more balanced division of voters

## Y axis: discrepancy $\beta$

the higher the more conflict each pair generates


Sushi


Skate


|  | 2017 |  |  | 2022 |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| MaxSwap | Far-left | ¢ | Far-right | Far-left | $\boldsymbol{j} \rightarrow \boldsymbol{j}$ Far-right |
| MaxNash | Socialist | ¢ | Far-right | Left | j $\sim$ Far-right |
| MaxSum | Socialist | ¢ | Far-right | Far-left | $j \rightarrow j$ Far-right |
| 2-MaxPolar | Far-left | i -1 | Far-right | Far-left | $\boldsymbol{j} \rightarrow \boldsymbol{j}$ Far-right |
| Borda | Left | $i \rightarrow$ | Liberal | Left | $\boldsymbol{j} \boldsymbol{j} \boldsymbol{j}$ Green |
| CC | Left |  | Conservative | Green | $\boldsymbol{j} \rightarrow \boldsymbol{j}$ Far-right |

Figure Skating


Political


Sushis

recap

## Axioms

## Selection Rules

Theoretical Validation

## Experimental Validation

## More than two candidates?

## Approval setting?

## *Relation to full-election measures?

*Categorizing elections?

## Thank you!

