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# NP-hard metrics for evaluating PB election rules

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# Preliminaries

- $E = (P, N, b, cost)$  – election
- $P = \{p_1, \dots, p_m\}$  – projects
- $N = \{1, 2, \dots, n\}$  – voters
- $b \in \mathbb{N}$  – budget
- $cost : P \rightarrow \mathbb{N}$
- $s_i(p) \in \mathbb{N}$  – score

## Score utilities

$$u_i^{sc}(W) = \sum_{p \in W} s_i(p)$$

## Cost utilities

$$u_i^{cost}(W) = \sum_{p \in W} s_i(p) cost(p)$$

# Voting rules

voting rule – function  $E \rightarrow P(P)$

$W \in P(P)$  – outcome, s.t.  $cost(W) \leq b$

Outcome  $W$  is **exhaustive** if for each project  $c \in C \setminus W$  we have that  $cost(W) + cost(c) > b$

A voting rule is **exhaustive** if it always returns an exhaustive outcome.

Two popular voting rules are:

- Utilitarian Greedy,
- Method of Equal Shares.

# Utilitarian Greedy

1. We start with an empty outcome  $W = \emptyset$ .
2. Repeatedly select a project  $p$  maximising the ratio  $\sum_{i \in N} \frac{u_i(p)}{\text{cost}(p)}$ .
3. If  $\text{cost}(W) + \text{cost}(p) \leq b$  then add project  $p$  to  $W$ ; otherwise remove the project from consideration.

This rule aims at maximizing the total utility of voters.

It is **optimal up to one project** – for each outcome  $W$  returned by UG there exists  $p \notin W$  s.t.:

$$\sum_{i \in N} u_i(W \cup \{p\}) \geq \max_{W': \text{cost}(W') \leq b} \sum_{i \in N} u_i(W').$$

# Basic Metrics for Fairness and Efficiency

**Average utility:**

$$\frac{1}{n} \sum_{i \in N} u_i(W)$$

**Dominance margin** of  $R_1$  over  $R_2$ :

the fraction of voters who enjoy strictly higher utility from the outcome of  $R_1$  than from  $R_2$ .

Similar ones: **improvement margin** and **exclusion ratio**.

Let's say, for a given outcome  $W$ , voter  $i$ 's share is:

$$share_i(W) = \sum_{p \in W} \frac{s_i(p)}{\sum_{j \in N} s_j(p)} \cdot cost(p)$$

**Power inequality:**

$$\frac{1}{n} \cdot \sum_{i \in N} \left| share_i(W) - \frac{b}{n} \right| \cdot \frac{n}{b}$$

# Basic Metrics for Fairness and Efficiency

Let  $\mathcal{D} = \{D_1, \dots, D_t\}$  be the set of districts (a partition of  $N$ ). The **dispersion of the budget allocation**:

$$\frac{1}{|\mathcal{D}|} \cdot \sum_{D \in \mathcal{D}} \frac{|\sum_{i \in D} \text{share}_i(W) - |D|/n \cdot b|}{|D|/n \cdot b}$$

City	Add1U, C	Util. G, D	Util. G, C
Czestochowa	0.23	0.28	0.39
Gdansk	0.27	0.33	0.46
Katowice	0.19	0.26	0.51
Krakow	0.08	0.24	0.23
Warsaw	0.20	0.41	0.41
Wroclaw	0.15	0.26	0.22
Zabrze	0.38	1.24	0.41

Let  $W_{sc}$  and  $W_{appr}$  be the outcomes of a given voting rule for the original and the approval elections, respectively. The **robustness ratio**:

$$\frac{\text{cost}(W_{appr} \cap W_{sc})}{\text{cost}(W_{sc})}$$

City	Add1U, C	Util. G, D	Util. G, C
Czestochowa	0.80	0.35	0.39
Gdansk	0.87	0.26	0.39
Katowice	0.83	0.56	0.42
Krakow	0.78	0.52	0.41

# Budget distribution among Categories

For each project  $p$ , denote by  $tags(p)$  the tags assigned to  $p$ . For each tag  $t$ , we can compute **vote share**:

$$\frac{1}{n} \sum_{i \in N} \sum_{p \in A_i: t \in tags(p)} \frac{1}{|A_i| \cdot |tags(p)|}$$

and also for an outcome  $W$  we define **spending share**:

$$\frac{1}{cost(W)} \sum_{p \in W: t \in tags(p)} \frac{cost(p)}{|tags(p)|}$$

With these two, we can compute  $l_2$  distance between the vectors of vote shares and the spending shares of all the tags.

	Public space	Sport	Transit	Education
Vote share	20%	16%	5%	33%
Equal Shares	24%	19%	11%	26%
Utilitarian Greedy	30%	26%	15%	5%

# The Core

For an election  $(N, C, b, cost)$  we define  
**extended justified representation:**

$$\forall E \forall l \in \mathbb{N} \forall S \subseteq N: S \text{ is } l\text{-cohesive} \Rightarrow \exists i \in S |A(i) \cap R(E)| \geq l$$

Committee  $W$  is in the **core** if

$$\forall S \subseteq N \forall T \subseteq C \left( \frac{|S|}{|N|} \geq \frac{cost(T)}{b} \Rightarrow \exists i \in S u_i(W) \geq u_i(T) \right)$$

Is the core always non-empty?

What we know:

- Core can be empty in elections with ranking-based preferences and in PB with additive utilities
- $Core \implies EJR$



# Pabutools and Master's Thesis

- Pabutools – Python library with implementation of some rules for PB
- Implement metrics for evaluating elections using linear programming
- Research whether the core can be computed fast in some cases (maybe with additional assumptions / constraints)
- Possibly research other metrics such Fractional Core, Pareto-Optimality, etc.

**Thanks for your attention**